

Teaching

Mo: 9:45 - 11:15
Fr 11:30 - 15:00

- Today:
- 1) Move on plane curves
 - 2) Approach to schemes via functor of points
 - 3) Fiber products, products
- Next Morphisms of schemes, properties.

1) Rational parametrization of conics quadratic curves in \mathbb{P}^2

Again $Q =$ quadratic form in x, y, z over alg. closed field can be transformed to

$$(x^2 + y^2 - z^2)$$

$$x^2 + (y-z)(y+z)$$

↳ change of variables

$$x^2 + yz$$

Sometimes it is better to make quadratic forms not $\approx \sum_{i=1}^n X_i^2$, but

$$X_1^2 + X_2 X_3 + X_4 X_5 + \dots + X_{2k} X_{2k+1} \quad (n \text{ odd})$$

$$X_1 X_2 + X_3 X_4 + \dots \quad (n \text{ even}).$$

Given a conic, a point on a conic, and a line not containing the point

ie. on $x^2 + yz = 0$ $(0, 1, 0) = p_0$

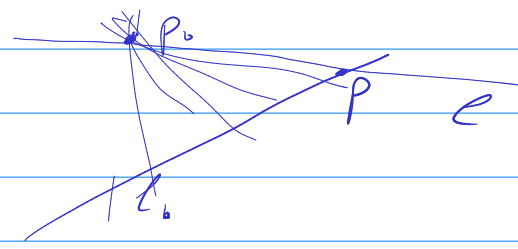
line $\{(s, 0, t) \mid (s, t) \neq (0, 0)\} = l_0$ $\mathbb{P}^1 \subset \mathbb{P}^2$

Let $p = (s, 0, t) \in l_0$

Then l is span of (p_0, p) , i.e.

l is parametrized by

$$p_0 \cdot u + p \cdot v \quad \{(vs, u, vt) \mid (u, v) \neq (0, 0)\} = l$$



$$L \cap Q: (vs)^2 + u \cdot vt = 0, \text{ solve for } u, v.$$

$$v(vs^2 + ut) = 0 \quad v=0 \text{ is always a solution}$$

corresponds to $(0, u, 0) \text{ or } (0, 0, 0)$

$vs^2 + ut$ gives another intersection point of Q and L " p_0

$$u = s^2, v = -t$$

plug in, $(vs, u, vt) = (-st, s^2, -t^2)$

So we obtained a map $\mathbb{P}^1 \rightarrow \mathbb{P}^2$ given by sending (s, t) to $(-st, s^2, -t^2)$. This is a bijection on Q . (To see this, we can cover Q by open sets, on each set define an inverse.)

$$Q \quad (x^2 + yz = 0) \quad x \neq 0 = U_x \quad \frac{y}{x} = \alpha \quad \frac{z}{x} = \beta$$

equation $2\beta + 1 = 0$

p_0

$$L = u(1, \alpha, \beta) + v(0, 1, 0)$$

$$(u, u\alpha + v, u\beta)$$

$$L \cap L_0$$

$$v = -u\alpha$$

$$(u, 0, u\beta) = p$$

$$(1, \beta) \in \mathbb{P}^1$$

on $U_x \quad (\alpha, \beta) \rightarrow (1, \beta) \in \mathbb{P}^1$

on $U_y \quad \frac{x}{y} = \gamma \quad \frac{z}{y} = \delta \quad \text{equation } \gamma^2 + \delta = 0$

compute the map on U_y

$$L = u(\gamma, 1, \delta) + v(0, 1, 0) \quad L \cap L_0: v = -u$$

$$U_y \ni (\gamma, \delta) \rightarrow (\gamma, \delta) \in \mathbb{P}^1$$

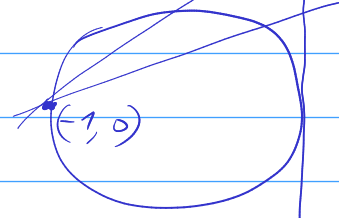
on $U_z \quad \frac{x}{z} = \epsilon, \frac{y}{z} = \zeta \quad \text{eq. } \epsilon^2 + \zeta = 0$

$$L = u(\epsilon, \zeta, 1) + v(0, 1, 0) \quad (\epsilon, 1)$$

$$\epsilon, \zeta \in U_z \text{ go to } (\epsilon, 1) \in \mathbb{P}^1$$

Look again $(s, t) \rightarrow (-st, s^2, -t^2)$

More traditional
parametrize circle
by points on line



$$x^2 + y^2 = 1 \quad \left(\frac{2s}{s^2+1} \right)^2 + \left(\frac{s^2-1}{s^2+1} \right)^2 = 1 \quad \mathbb{Q}$$

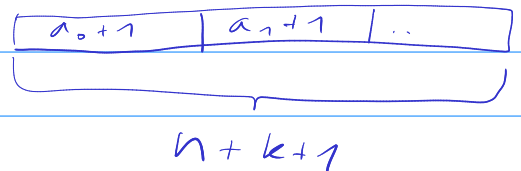
so $s \rightarrow \left(\frac{2s}{s^2+1}, \frac{s^2-1}{s^2+1} \right)$ parametrizes the circle.
 ← be same parametrization in different coordinates

used to classify pythagorean triples
(good example how alg. geometry helps number theory)

Remark changing coordinates we can
consider $(s, t) \rightarrow (s^2, st, t^2)$ This parametrizes
 $x \quad y \quad z \quad xz = y^2$

This is example of Veronese embedding.

More generally, take \mathbb{P}^k , coordinates (x_0, \dots, x_k) .
list all monomials of degree n . There are
gives a map $\mathbb{P}^k \rightarrow \mathbb{P}^{\binom{n+k}{k}-1} = N$
 $x_0^{a_0} \dots x_k^{a_k}$



This is an embedding
(so \mathbb{P}^k is realized as
a closed subscheme of \mathbb{P}^N).

there are $n+k$ positions
to put k vertical lines
so $\# = \binom{n+k}{k}$.

More generally

If R is a graded ring, let $n \in \mathbb{Z}_{\geq 0}$, let
 $R^{(n)} = \bigoplus_{k=0}^{\infty} R_{n+k}$. Then $\text{Proj}(R^{(n)}) \cong \text{Proj}(R)$.
why? Let U_{\pm} be a basic open set of $\text{Proj}(R)$

let f of degree k . Then $f^n \in R^{(n)}$
 clearly $\left[(f^n)^{-1} R^{(n)} \right]_0 = \left[f^{-1} R \right]_0$, so $U_f \cong \left[U_{f^n}^{R^{(n)}} \right]$
 $\left[(f^n)^{-1} R \right]_0$ corresponding basic open of $\text{Proj } R^{(n)}$.

So for $R = k[x_0, \dots, x_n]$ $R^{(n)}$ is generated by all monomials of degree n , so it is a quotient of the polynomial ring in $N+1$ variables by some ideal, so \mathbb{P}^n is isomorphic to a closed subscheme of \mathbb{P}^N .

Terminology: a scheme is projective if it is a closed subscheme of a projective space (equivalently, it is Proj (quotient of $k[x_0, \dots, x_n]$), equivalently, it is Proj (finitely generated graded ring satisfying $R_0 = k$).

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Ex

weights proj space \subset projective space

$\text{Proj } k[x_0, \dots, x_n]$

$\deg x_i = d_i$

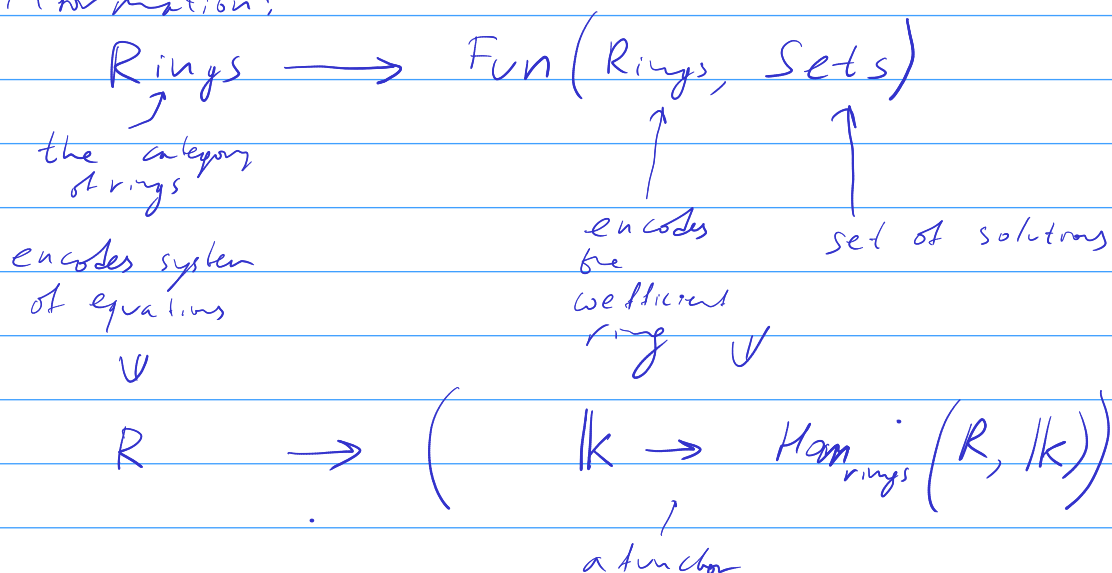
Let $L = \text{l.c.m.}(d_0, \dots, d_n)$, list all monomials of degree L . If there are N of them obtain embedding into \mathbb{P}^{N-1} .

May be L doesn't work, but some multiple of L does.

(want every monomial of degree kL be divisible by some monomial of degree L). good exercise...

Functor of points.

Recall if we have a system of equations, we should study solutions over arbitrary rings, Yoneda lemma guarantees that we don't lose information:



This is a contravariant functor.

(if $\varphi: R_1 \rightarrow R_2$) then $\forall K$ we have a map of sets $\text{Hom}(R_2, K) \rightarrow \text{Hom}(R_1, K)$.

for each R the functor $K \rightarrow \text{Hom}(R, K)$ is called the functor of points.

We don't lose information means that by Yoneda lemma the category of rings embeds into the category of functors $\text{Rings} \rightarrow \text{Sets}$.

Embeds := bijection on Hom

injection : faithful

surjection : full

What about schemes?

We have a category of schemes: Objects are schemes,
Denoted by Sch . Morphisms are morphisms

We have a full subcategory of schemes.
 $AffSch$ of affine schemes.

Objects of $AffSch$ are schemes of the form $Spec(R)$ for rings R .

We know that $\forall R_1, R_2$

$$\boxed{\text{Mor}(Spec(R_1), Spec(R_2)) = \text{Hom}(R_2, R_1)}$$

So $AffSch$ is isomorphic to $\text{Rings}_{\mathbb{A}^1}$
morphisms go in the opposite direction.

The functor of points can be interpreted as

$$\begin{array}{ccc} AffSch & \longrightarrow & (\text{Rings} \rightarrow \text{Sets}) \\ \downarrow & & \downarrow \\ Spec(R) & & \mathbb{k} \rightarrow \text{Hom}(R, \mathbb{k}) \\ X & \longrightarrow & (\mathbb{k} \rightarrow \text{Mor}(Spec(\mathbb{k}), X)) \end{array}$$

We generalize this to arbitrary schemes by the same formula

$$\begin{array}{ccc} Sch & \longrightarrow & (\text{Rings} \rightarrow \text{Sets}) \\ X & \longrightarrow & (\mathbb{k} \xrightarrow{F_X} \text{Mor}(Spec(\mathbb{k}), X)) \end{array}$$

Variation of this Schemes over a field, instead of Mor take Mor over that field.

Theorem This functor is an embedding of Sch into $\text{Fun}(\text{Rings}, \text{Sets})$.

Sketch of proof

We need to understand what are morphisms between two arbitrary schemes X_1, X_2 in terms of morphisms between affine schemes.

X_1 is a union of open affine schemes $\{U_\lambda\}_{\lambda \in \Lambda}$
 $\text{Mor}(X_1, X_2) = \left\{ (f_\lambda : U_\lambda \rightarrow X_2)_{\lambda \in \Lambda} \mid \forall \lambda, \mu \in \Lambda \right.$
 Let $U_\lambda = \text{Spec } R_\lambda$
 $f_\lambda|_{U_\lambda \cap U_\mu} = f_\mu|_{U_\lambda \cap U_\mu}$
 $f_\lambda \in F_{X_2}(R_\lambda)$

$U_\lambda \cap U_\mu$ does not have to be affine, however can be covered by affines
 Let's rewrite

Exercise intersection of two open affines doesn't have to be affine.

$$f_\lambda|_{U_\lambda \cap U_\mu} = f_\mu|_{U_\lambda \cap U_\mu}$$

$$\underbrace{\mathbb{A}^2} \cup \underbrace{\mathbb{A}^2}$$

$$U = \mathbb{A}^2 \setminus \{(0,0)\}$$

as follows: for any affine scheme V and any commutative diagram

$$\begin{array}{ccc} V \rightarrow U_\lambda & & \text{we want} \\ \downarrow & & \downarrow \\ U_\mu \rightarrow X_1 & & \begin{array}{ccc} V \rightarrow U_\lambda & & \\ \downarrow & & \downarrow f_\lambda \\ U_\mu \rightarrow X_2 & & \\ \uparrow f_\mu & & \end{array} \end{array}$$

\forall ring R and homomorphisms $\alpha: R_\lambda \rightarrow R, \beta: R_\mu \rightarrow R$ whenever

$$\text{such that } \alpha(U_\lambda \rightarrow X_1) = \beta(U_\mu \rightarrow X_1)$$

to be commutative

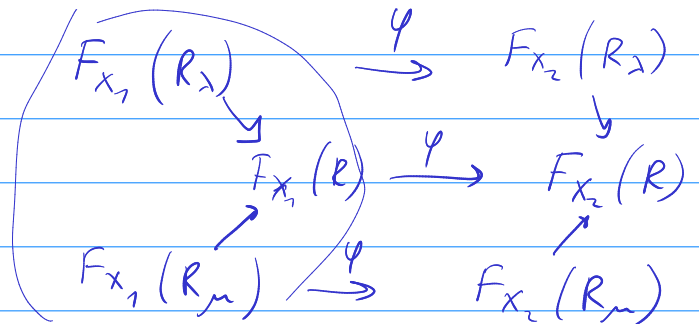
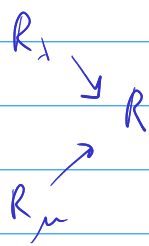
$$\underbrace{F_{X_2}(R_\lambda)}_{\uparrow} \\ F_{X_2}(R)$$

also $\alpha(f_\lambda) = \beta(f_\mu)$ holds.

$$\text{Mor}(X_1, X_2) \xrightarrow{\quad \varphi \quad} \text{Mor}(F_{X_1}, F_{X_2})$$

1) why injective? because each map $X_1 \rightarrow X_2$ is uniquely determined by the collection $\{t_\lambda\}$, where t_λ is nothing but the value of φ on $U_\lambda \rightarrow X_1$.
 $t_\lambda = \varphi(U_\lambda \rightarrow X_1)$ \uparrow
 $F_{X_1}(R_\lambda)$

2) why surjective? because given a collection t_λ , it defines a morphism $X_1 \rightarrow X_2$ when some conditions are satisfied.



Commutative diagram because φ is a natural transformation.

the condition is clear from the commutativity.

Theorem is proved.

Remark In categorical language X_1 is a colimit of all of its open affine subschemes.

Basically we used that morphisms $\text{Mor}(-, X_2)$ commutes with colimits.

When $F = F_X$ is given we say X represents the functor F .

So we have a way to construct schemes by

- 1) Construct a functor
- 2) Prove that the functor is representable.

Examples

Fiber product

Setup

X_1, X_2, Y schemes with
morphisms $\pi_1: X_1 \rightarrow Y$ and $\pi_2: X_2 \rightarrow Y$. Consider the
functor

$$F(R) = \left\{ (t_1 \in F_{X_1}(R), t_2 \in F_{X_2}(R)) \mid \pi_1(t_1) = \pi_2(t_2) \right\} \cong F_Y(R)$$

Theorem F is representable by a scheme called the fiber product $X_1 \times_Y X_2$.

Example $Y = \text{Spec } \mathbb{Z}$ (Spec k if we are over a field k)
Then this is called the product for any scheme X there is a unique map to Y , so $\pi_1(t_1) = \pi_2(t_2)$ is automatic.

Example $A^k \times A^n = A^{k+n}$
 \forall ring R $F_{A^k \times A^n}(R) = \underbrace{R \times \dots \times R}_k$

Example $G_m = A^1 \setminus \{0\} = \text{Spec } \mathbb{Z}[t, t^{-1}]$ ($= \text{Spec } k[t, t^{-1}]$ if we are over k)

$$F_{G_m}(R) = \{x \in R \mid x^{-1} \text{ exists}\} = R^\times$$

Note that R^\times is a commutative group, so in particular we have a map $R^\times \times R^\times \rightarrow R^\times$ (multiplication)

So automatically we obtain a morphism of schemes

$$G_m \times G_m \rightarrow G_m \quad (\text{explicitly this corresponds to a homomorphism of rings } k[t, t^{-1}] \rightarrow k[u, u^{-1}, v, v^{-1}])$$

$t \mapsto uv$

G_m is called a group scheme

$(\forall R \quad F_x(R) \text{ is a group})$

Proof of Theorem

$$X_1 \xrightarrow{\pi_1} Y \xleftarrow{\pi_2} X_2$$

Reduction 1 assume $U \subset Y \quad V_1 \subset \pi_1^{-1}(U)$

$V_2 \subset \pi_2^{-1}(U)$ are affine open.

$$U = \text{Spec}(R) \quad V_1 = \text{Spec}(R_1) \quad V_2 = \text{Spec}(R_2)$$

What is $V_1 \times_U V_2$? $= \text{Spec}(R_1 \otimes_R R_2)$

given rings R_1, R_2 , homomorphisms ψ

$$f_i: R \rightarrow R_i$$

$R_1 \otimes_R R_2$ = generated by symbols $x \otimes 1, 1 \otimes y$
 $x \in R_1, y \in R_2$, relations

$$(x_1 \otimes 1)(x_2 \otimes 1) = x_1 x_2 \otimes 1 \quad x_1 \otimes 1 + x_2 \otimes 1 = (x_1 + x_2) \otimes 1$$

similarly for $1 \otimes y_1, 1 \otimes y_2$

$$f_1(z) \otimes 1 = 1 \otimes f_2(z) \quad \forall z \in R.$$

(= tensor product as abelian groups).
because

$$x \otimes y := (x \otimes 1) \cdot (1 \otimes y) \quad \text{additively generate } R_1 \otimes_R R_2.$$

Example

$$k[x_1, \dots, x_n] \otimes_k k[y_1, \dots, y_m] = k[x_1, \dots, x_n, y_1, \dots, y_m].$$

$$R/I \otimes R' = (R \otimes R') / (I \otimes 1)$$

In general $X_1 \times_Y X_2$ is obtained by gluing the $U_1 \times_U U_2$.

Question is $\mathbb{P}^n \times \mathbb{P}^m$ projective?